

Post-Quantum Signatures in DNSSEC via Request-Based Fragmentation

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WHAT ARE WE TALKING ABOUT TODAY?

Background

- What is DNS?
- What (security) problems does DNS have?
 - How are these problems addressed?
 - Are these solutions future proofed?

PQC, DNSSEC, & Request-based Fragmentation

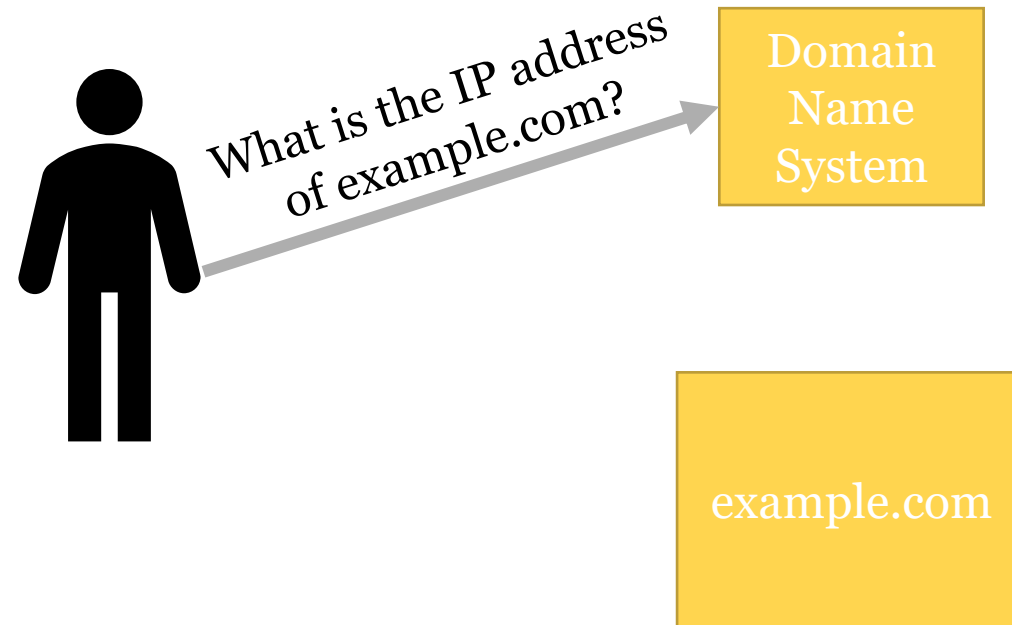
- Design
- Results & Evaluation
- Next Steps

WHAT IS DNS?

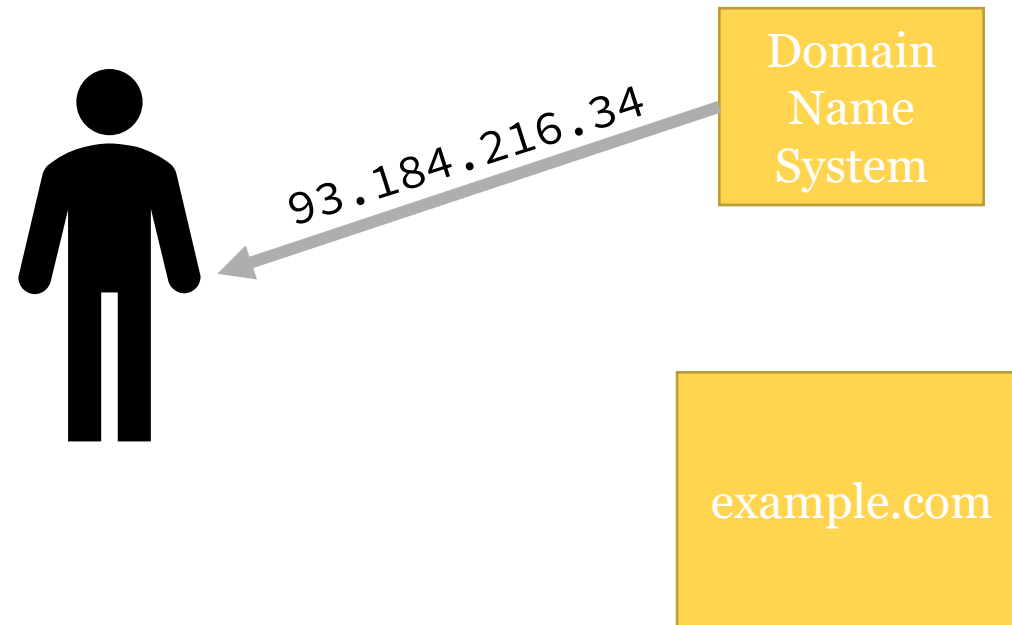
- The internet uses IP addresses to determine where to send messages
- IP addresses are difficult for people to remember!
- The Domain Name System is responsible to translating something easy for a human to remember into IP addresses

example.com -> 93.184.216.34

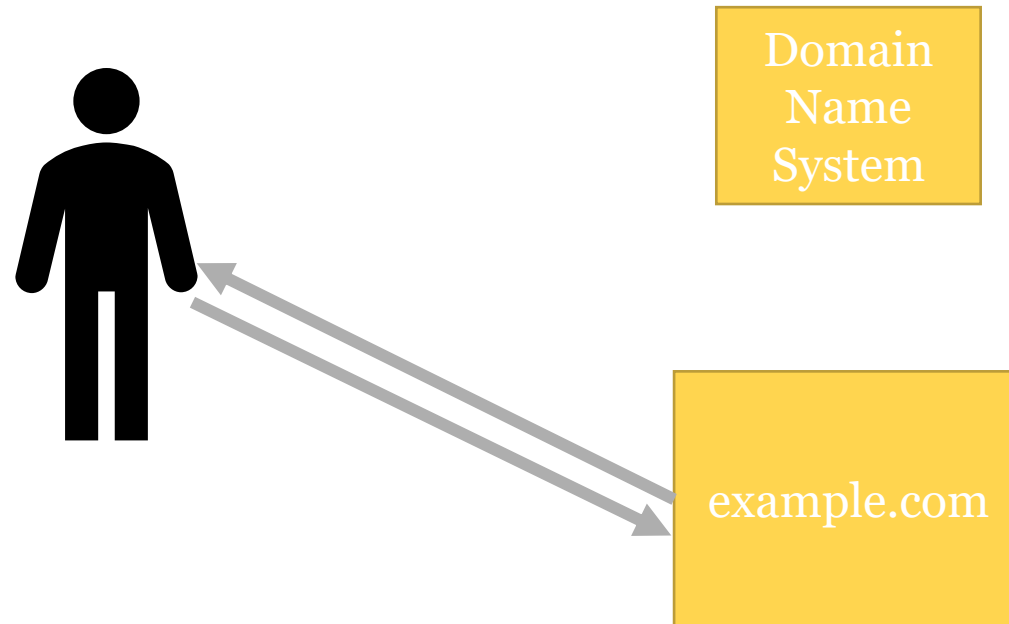
WHAT IS DNS?



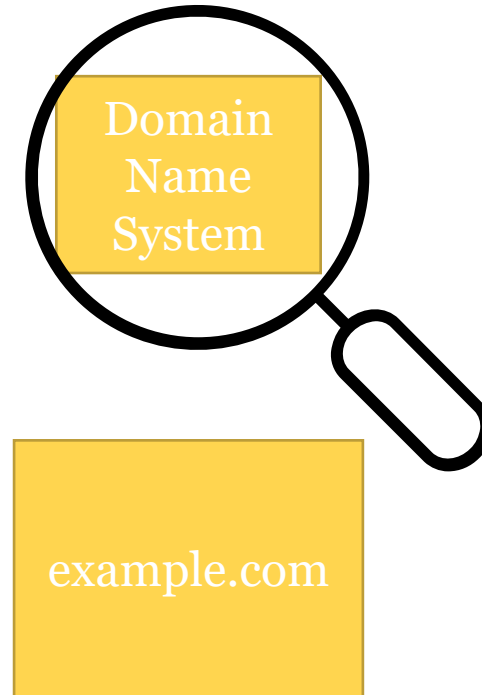
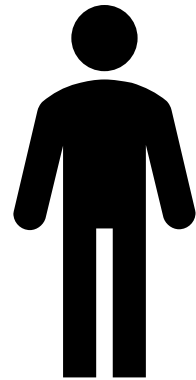
WHAT IS DNS?



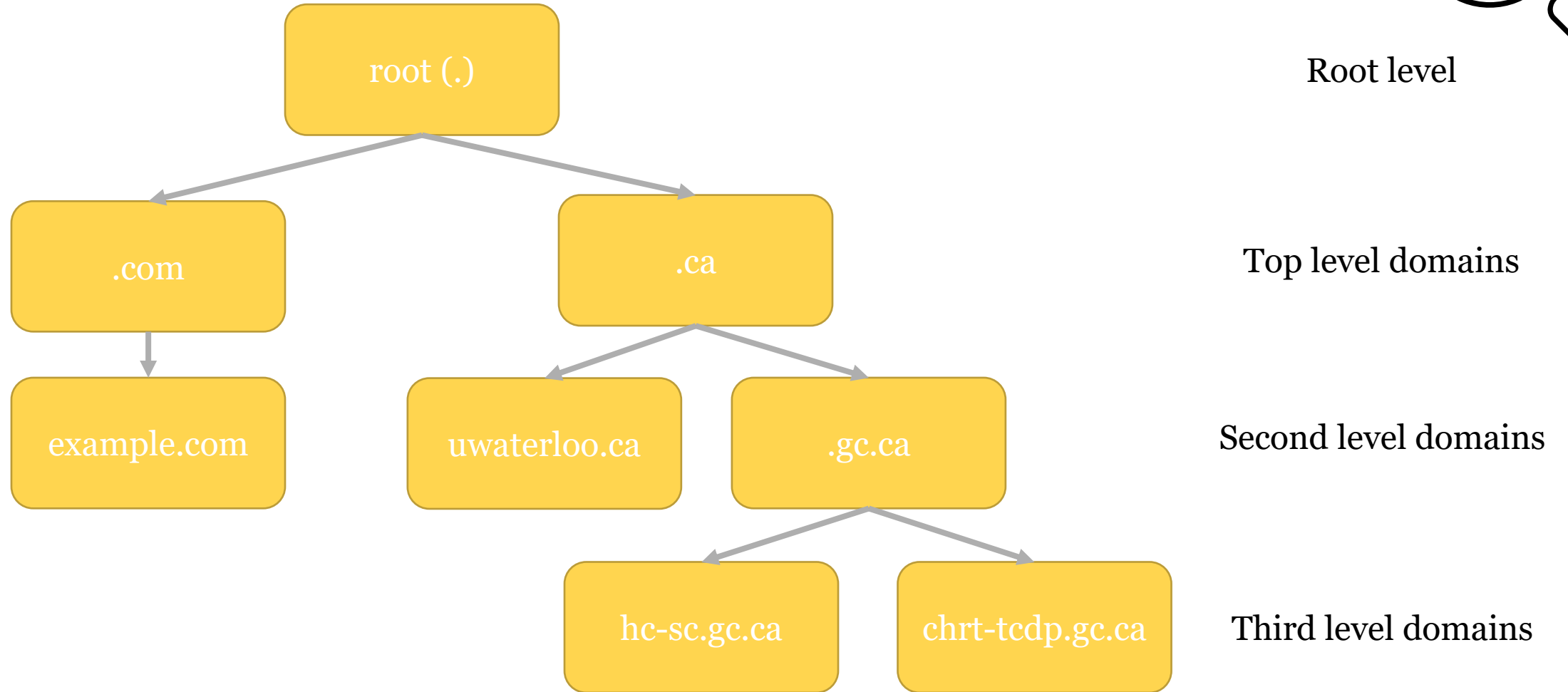
WHAT IS DNS?



WHAT IS DNS?



DNS IS BROKEN UP INTO ZONES



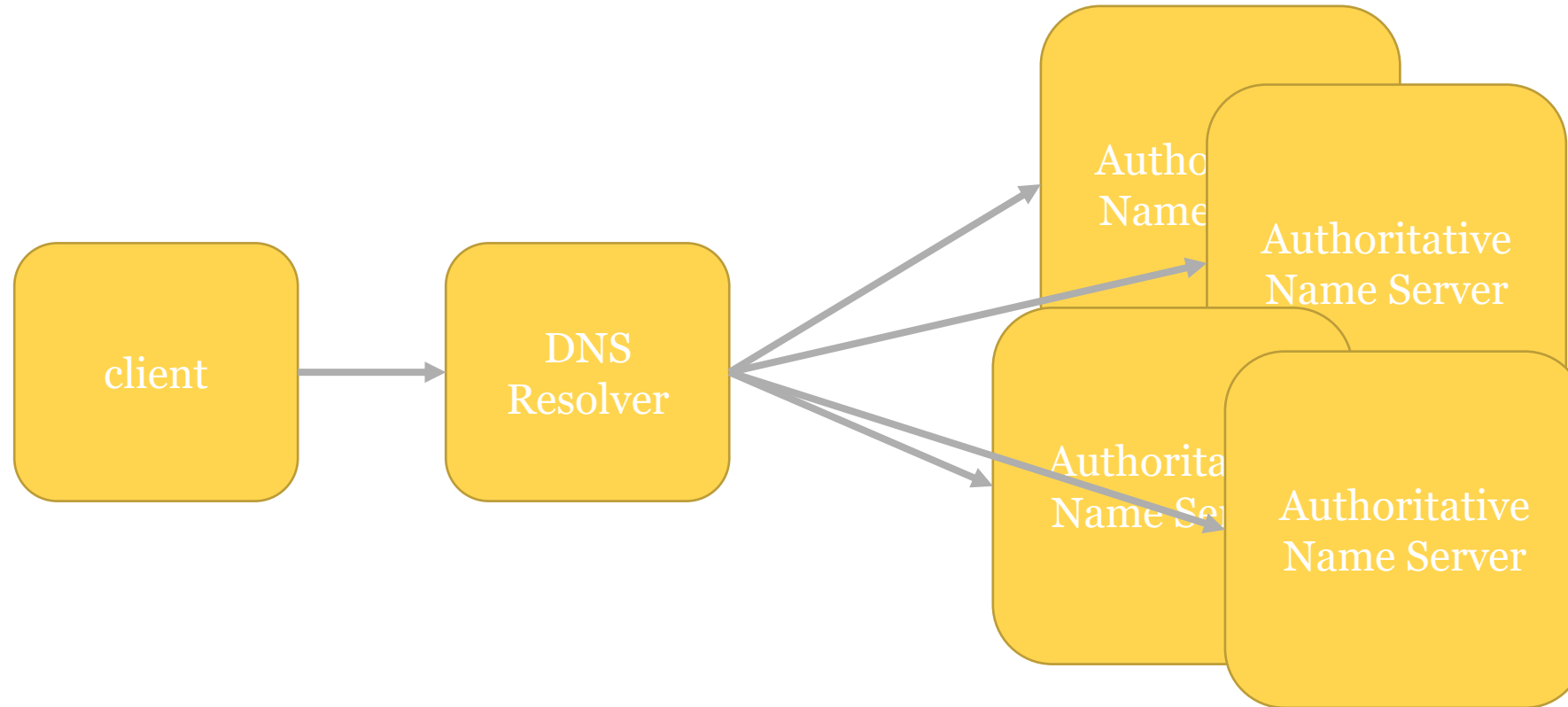
ZONES CONTAIN RESOURCE RECORDS



example.com. 57094 IN AAAA
example.com. 57047 IN A
example.com. 57094 IN NS
example.com. 57094 IN NS

2606:2800:220:1:248:1893:25c8:1946
93.184.216.34
b.iana-servers.net.
a.iana-servers.net.

CLIENTS RARELY QUERY DIRECTLY



PROBLEM WITH DNS



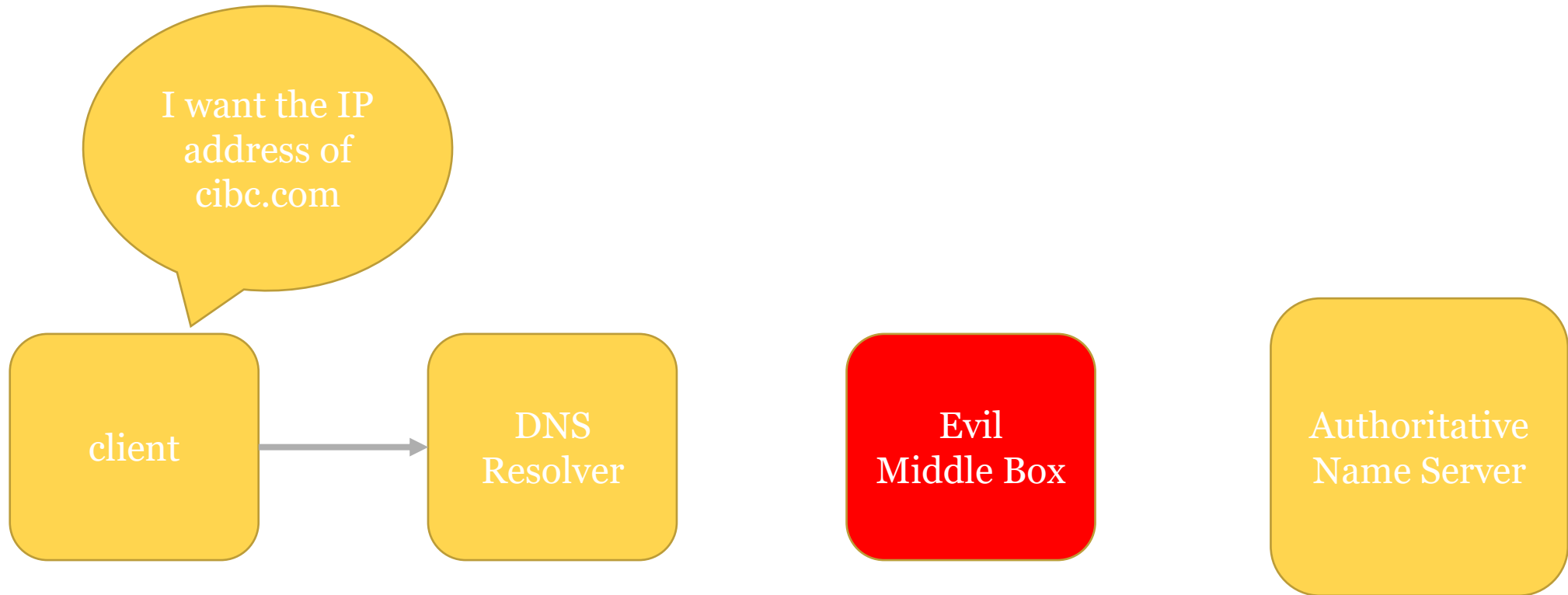
- Designed with no integrity protection



PROBLEM WITH DNS



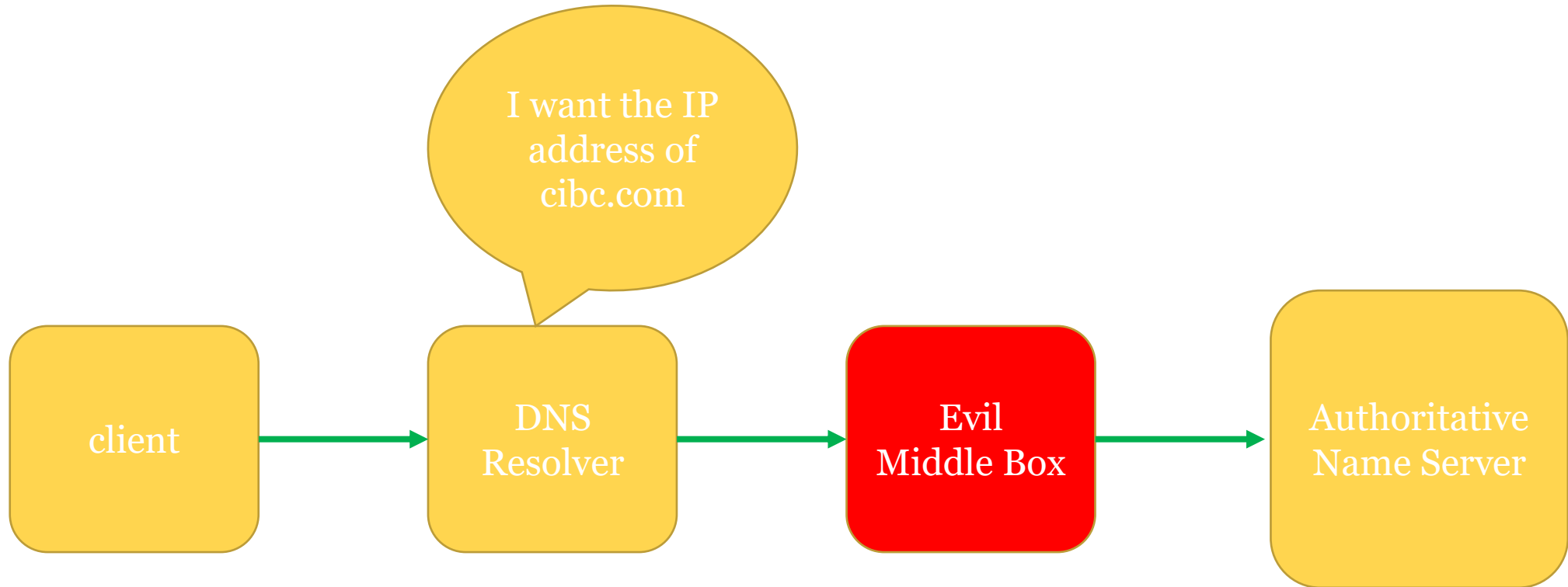
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PROBLEM WITH DNS



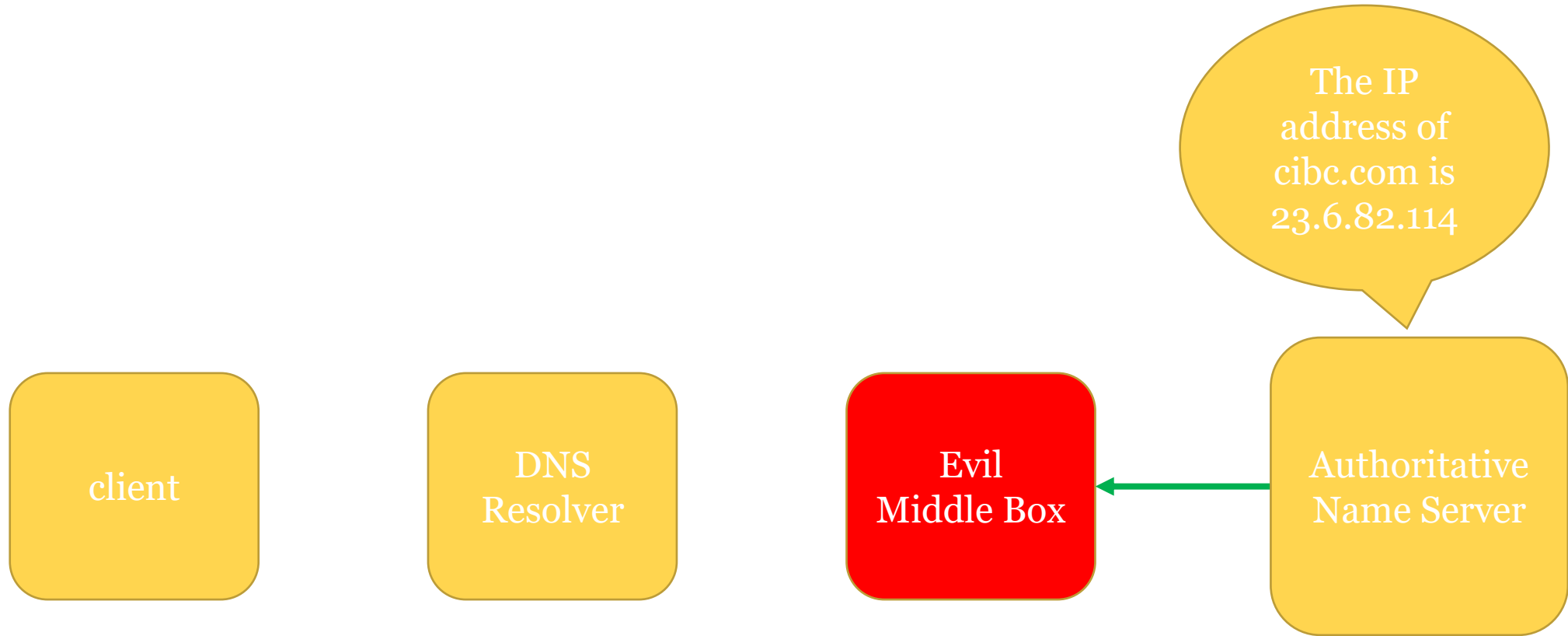
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PROBLEM WITH DNS



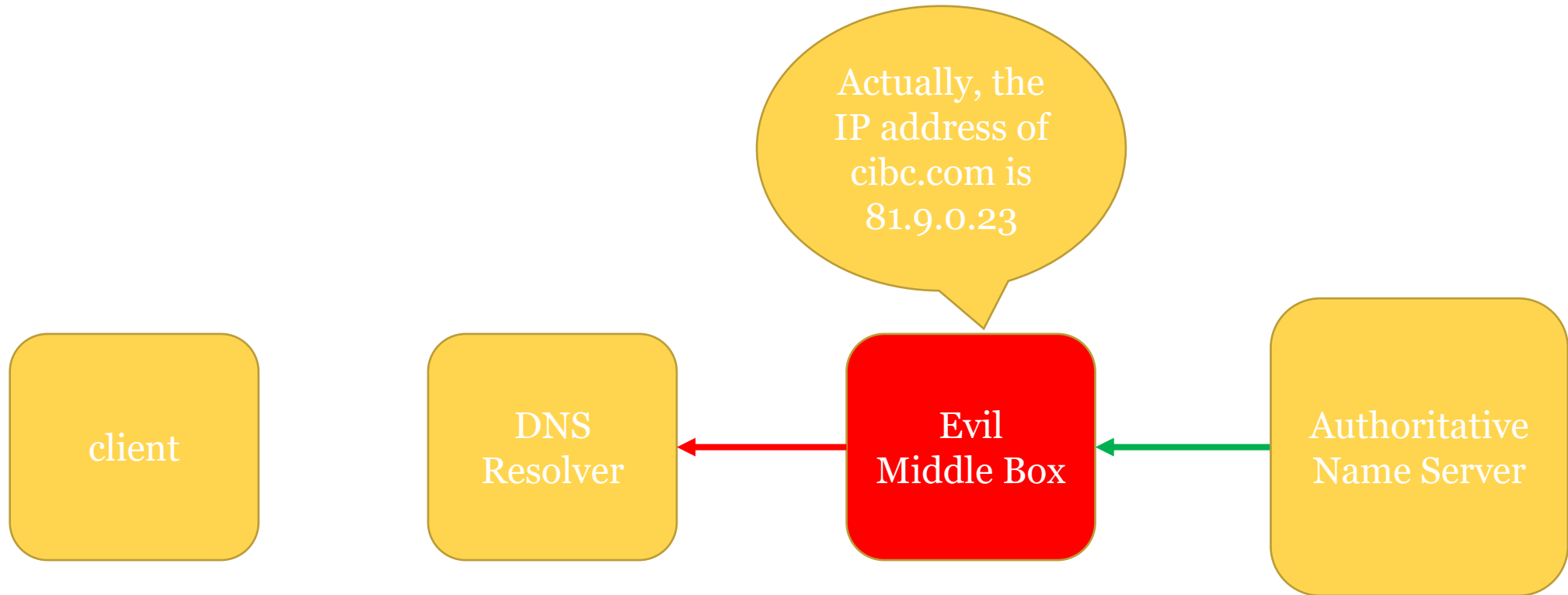
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PROBLEM WITH DNS



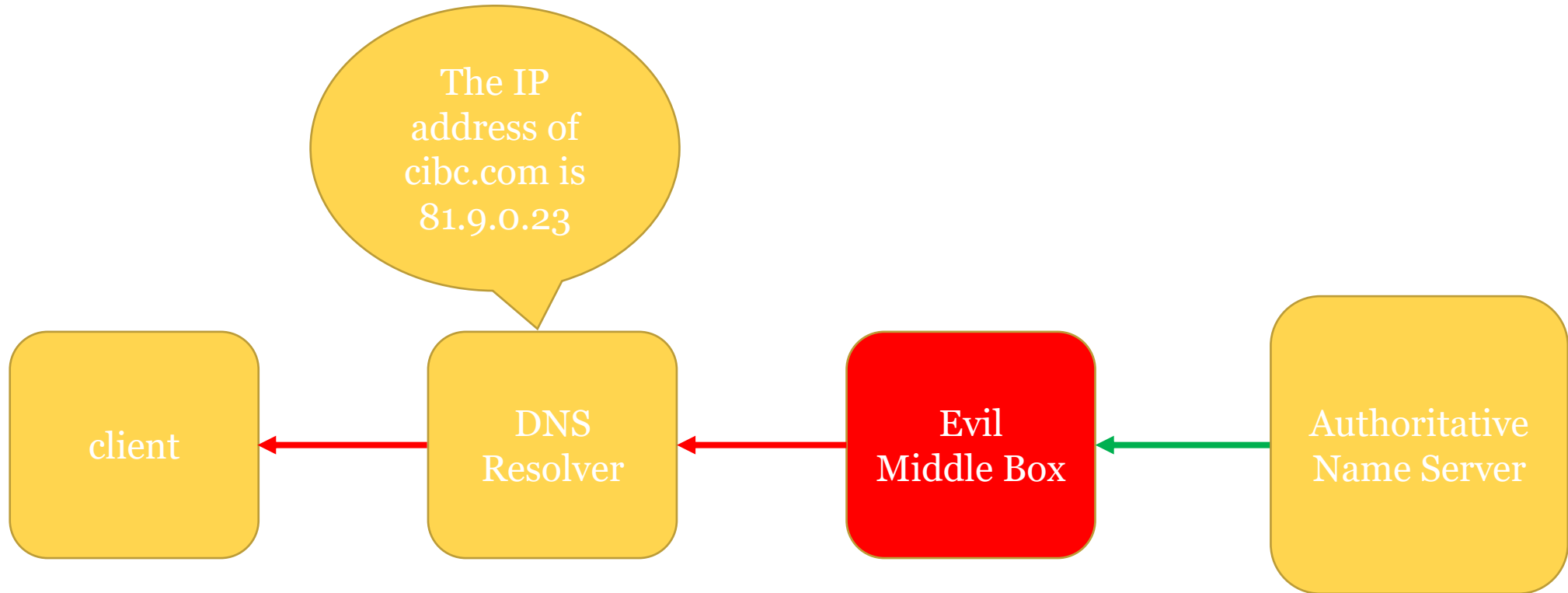
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PROBLEM WITH DNS



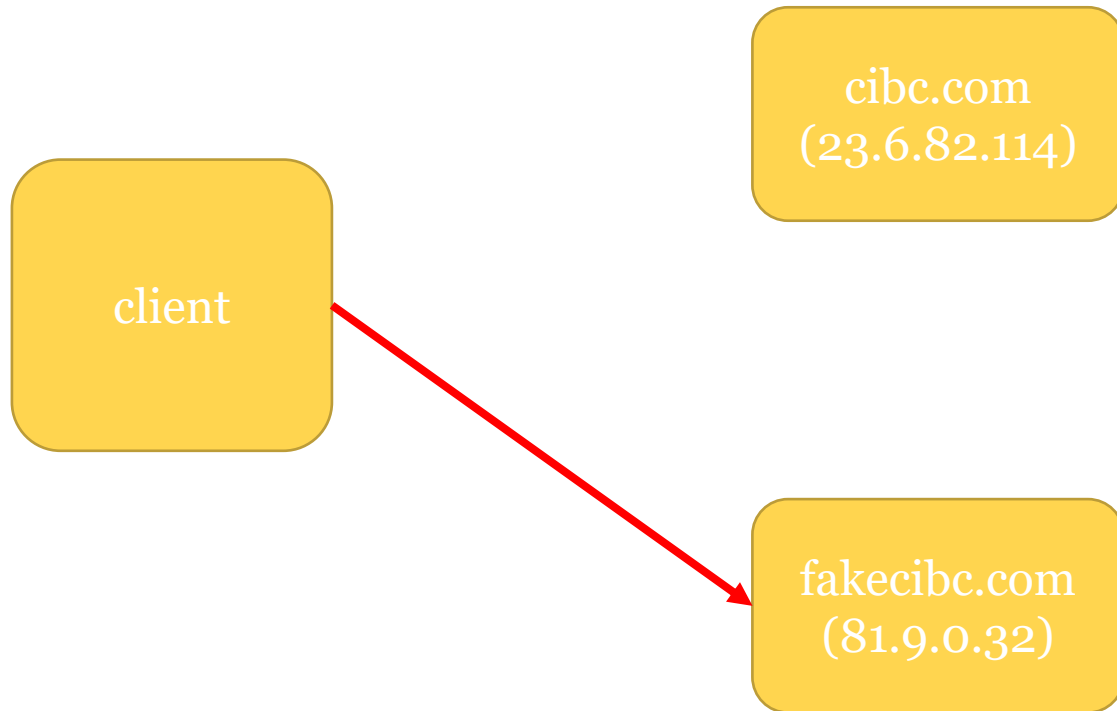
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PROBLEM WITH DNS



- Designed with no integrity protection



SOLVING DNS INTEGRITY

Use digital signatures to make sure the correct unmodified message is received and is from the correct entity!

- New records added to DNSSEC signed zone
- Sets of records (RRSets) are signed, rather than individual records
- Have two keys:
 - Key Signing Key
 - Zone Signing Key

SOLVING DNS INTEGRITY

example.com.	86379	IN	A	93.184.216.34
example.com.	86328	IN	RRSIG	A 8 2 86400 20221120061546 20221030132217 59208 example.com. rZvjehQxdT5pJ4cw+o1y/BYmLkBLuqzjFaEON9773Bhywt4qhKmME8DK oKD4yLjYJYFaqhUNCYb+iimCTdK+9+3UjJ35gRIDC3kuZ9hogtCoLBnt ltfgFwLQomdye8iH/FDDVKTm+CAz3UMfcwNzNahvg4BOnZo4HqnpZcWW pu4=
example.com.	73820	IN	NS	a.iana-servers.net.
example.com.	73820	IN	NS	b.iana-servers.net.
example.com.	86237	IN	RRSIG	NS 8 2 86400 20221122065049 20221101032317 59208 example.com. Uit3UXCeClM+iwVkc2lX8n5A1OoCD9mH8rsTSfcsjQaZD9y54q7bT2mM cGMaiyjj/sODGKLNvbFKLEgHgPoLNF4i+YzHvpct5MZD1c8JqnzYisOf xq+JQ4tLcsDmrnhJEinBVbiq/epEXso4I4GES+zyEgnz5TPErjTNRDzP 7CE=
example.com.	3600	IN	DNSKEY	257 3 8 AwEAAZoaqu1rJ6orJynrRfNpPmayJZoAx9Ic2/Rl9VQWLMHyjxxem3VU SoNUIFXERQbj0A9Ogp0zDM9YIccKLRd6LmWiDct7UJQxVdD+heb5Ec4q lqGmyX9MDabkvX2NvMwsUecbYBq8oXeTT9LRmCUt9KUt/WOi6DKECxoG /bWTykrXyBR8elD+SQY43OAVjlWrVltHxgp4/rhBCvRbmdflunaPIgu2 7eE2U4myDSL8a4AorB5uHG4PkOa9dIRs9y00M2mWf4lyPee7vi5few2 dbayHXmieGcaAHrx76NGAABeY393xjlmDNcUkF1gpNWUla4fWZbbaYQz A93mLdrng+M=
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example.com.	2694	IN	RRSIG	DNSKEY 8 2 3600 20221129010414 20221107233521 31406 example.com. nMEQXWFatPZd/fkGgi9TI4ZO2vokX+6zNNmZPSONweki1Vb25f+oISgH b1WEg84IzyUw+zzwmS2G4Jo8PvS8+rFfu9vprvPwKVMsGozBSyt3CCLS qa1DtY2oBMWXCzqHD1n16220AUMNGuvrta6ikmuGfXT/gXyK5isenUPn kSbGsbrgEQKPZZQU6H/9nLK2qttyBscCQmJ4ziIbsMyannBWgXtJgXhu 4AhiVAZIxCqII/ISNei3vOcl+h6C+RgjYsnoPD59HkpnC2H7TsaILNf7 uYtbCjzRKLhRzIwIS3ASbWccGJ3LXruZwUNdoE/XqrxacZXuwFrq+vtP RYAaPA==

SOLVING DNS INTEGRITY

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example.com.	73820	IN	NS	a.iana-servers.net.	
example.com.	73820	IN	NS	b.iana-servers.net.	
example.com.	86237	IN	RRSIG	NS 8 2 86400 20221122065049 20221101032317 59208 example.com. Uit3UXCeClM+iwVkc2lX8n5A1OoCD9mH8rsTSfcsjQaZD9y54q7bT2mM cGMaiyjj/sODGKLNvbFKLEgHgPoLNF4i+YzHvpct5MZD1c8JqnzYisOf xq+JQ4tLcsDmrnhJEinBVbiq/epEXso4I4GES+zyEgnz5TPErjTNRDzP 7CE=	
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example.com.	3600	IN	DNSKEY	256 3 8 AwEAAb1oJO+fCqdkxHtQYVB/tFPgJphc+VxjUYz+eVGf077zMxHKgce9 EwGBifFuKhjl2EAoVQPsWVX1vzuUmWri3OgsTBlITkdMz6VU4g94uO6T 9MIktokouOidIzvOqLR+O2LSXNhiYOIWA9s3Lxk5R2lrwd6vrRvT2CR1 GdZuUIKB	
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example.com.	73820	IN	NS	a.iana-servers.net.	} RRSet #2
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example.com.	73820	IN	NS	a.iana-servers.net.	} RRSet #2
example.com.	73820	IN	NS	b.iana-servers.net.	
example.com.	86237	IN	RRSIG	NS 8 2 86400 20221122065049 20221101032317 59208 example.com. Uit3UXCeClM+iwVkc2lX8n5A1OoCD9mH8rsTSfcsjQaZD9y54q7bT2mM cGMaiyjj/sODGKLNvbFKLEgHgPoLNF4i+YzHvpct5MZD1c8JqnzYisOf xq+JQ4tLcsDmrnhJEinBVbiq/epEXso4I4GES+zyEgnz5TPErjTNRDzP 7CE=	

example.com.	3600	IN	DNSKEY	257 3 8 AwEAAZoaqu1rJ6orJynrRfNpPmayJZoAx9lc2/Rl9VQWLMHyjxxem3VU SoNUIFXERQbj0A9Ogp0zDM9YIccKLRd6LmWiDct7UJQxVdD+heb5Ec4q lqGmyX9MDabkvX2NvMwsUecbYBq8oXeTT9LRmCUt9KUt/WO16DKECxoG /bWTykrXyBR8elD+SQY43OAVjlWrVltHxgp4/rhBCvRbmdflunaPIgu2 7eE2U4myDSL8a4AorB5uHG4PkOa9dIRs9y00M2mWf4lyPee7vi5few2 dbayHXmieGcaAHrx76NGAABeY393xjlmDNcUkF1gpNWUla4fWZbbaYQz A93mLdrng+M=
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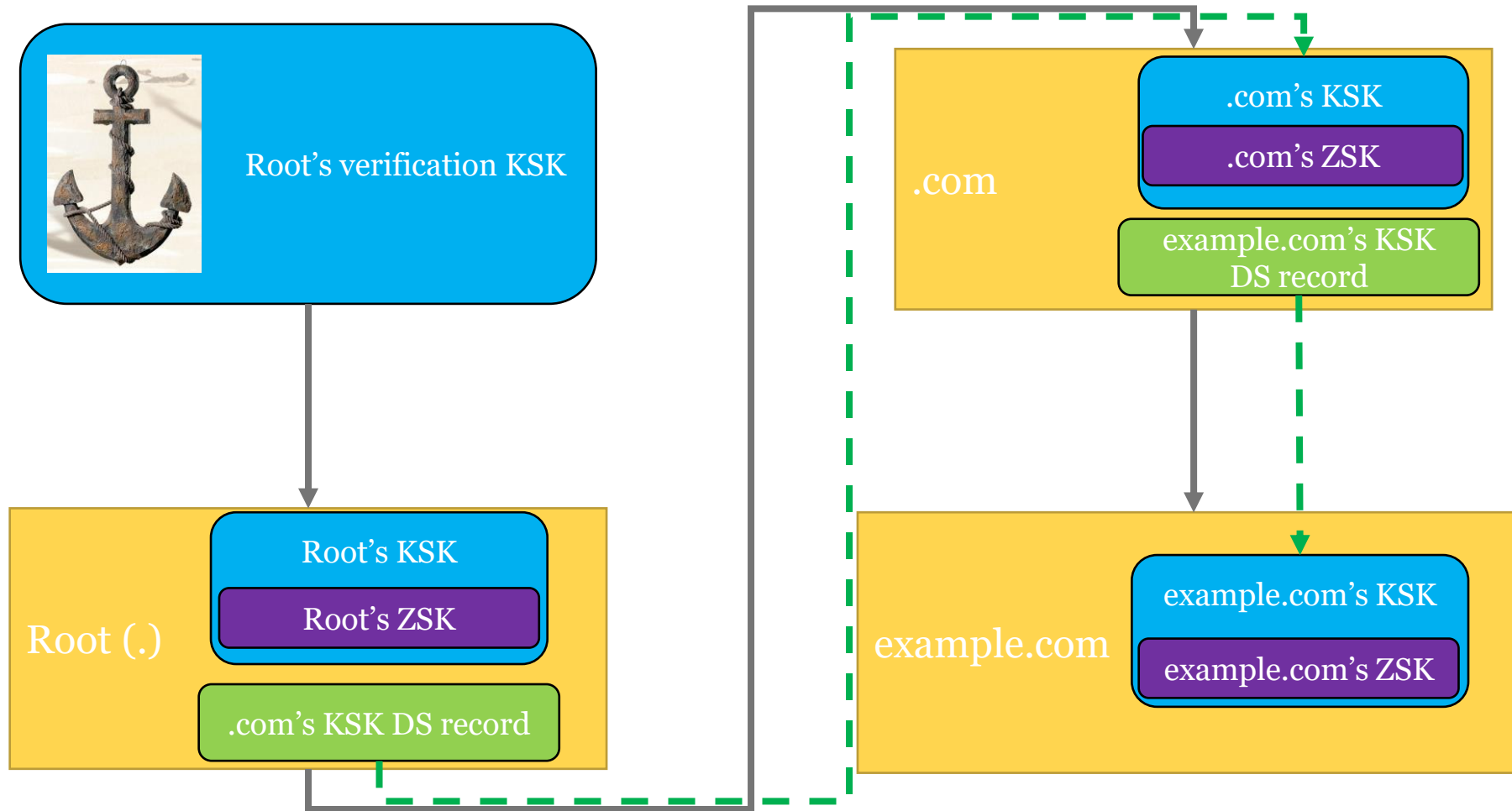
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example.com.	2694	IN	RRSIG	DNSKEY 8 2 3600 20221129010414 20221107233521 31406 example.com. nMEQXWFatPZd/fkGgi9TI4ZO2vokX+6zNNmZPSONweki1Vb25f+oISgH b1WEg84IzyUw+zzwmS2G4Jo8PvS8+rFfu9vprvPwKVMsGozBSyt3CCLS qa1DtY2oBMWXCzqHD1n1622oAUMNGuvrta6ikmuGfXT/gXyK5isenUPn kSbGsbrgEQKPZZQU6H/9nLK2qttyBscCQmJ4ziIbsMyannBWgXtJgXhu 4AhiVAZIxCqII/ISNei3vOcl+h6C+RgjYsnoPD59HkpnC2H7TsaILNf7 uYtbCjzRKLhRzIwIS3ASbWccGJ3LXruZwUNdoE/XqrxacZXuwFrq+vtP RYAaPA==

HOW DO WE MAINTAIN KEY INTEGRITY?

Construct a chain of trust!

- The root verification KSK must be manually configured on the machine making the request
- When the root ZSK is queried use the trust anchor to verify key and its signature
- Each zone's parent zone contains a "Delegate signer" (DS) record which is used to verify zone's KSK

HOW DO WE MAINTAIN KEY INTEGRITY?



POST-QUANTUM CRYPTO IN DNS

- We need to keep messages less than 1232 bytes
 - Prevents (fragile) UDP fragmentation
 - Avoids having to fallback to TCP (expensive)

Algorithm	Public Key Size	Signature Size
Falcon-512	897	666
Dilithium2	1,312	2,420
SPHINCS+-128s	32	7,856

HOW DO WE SEND LARGE PQC-SIGNED MESSAGES?

Idea: Let's move fragmentation from UDP to the DNS level

SENDING LARGE PQC DNS MESSAGES

What if we split large DNS messages into several smaller DNS messages (fragments)?

- Need to pass on extra info such as indicating support for this delivery method
- What if a resource record is **too** large to fit within the 1232 byte requirement?
- What if firewall is configured to only allow a single UDP based DNS response?

SENDING LARGE PQC DNS MESSAGES

What if we split large DNS messages into several non-fragmented UDP packets and blindly send them to the resolver to reassemble?

- How do we handle out of order delivery?
- What if a firewall is configured to only allow a single UDP packet for DNS messages?
- What if the resolver does not support this fragmentation method?
 - ICMP Flooding (Very bad!)

INTRODUCING **ARRF: A RESOURCE RECORD FRAGMENTATION MECHANISM**

ARRF

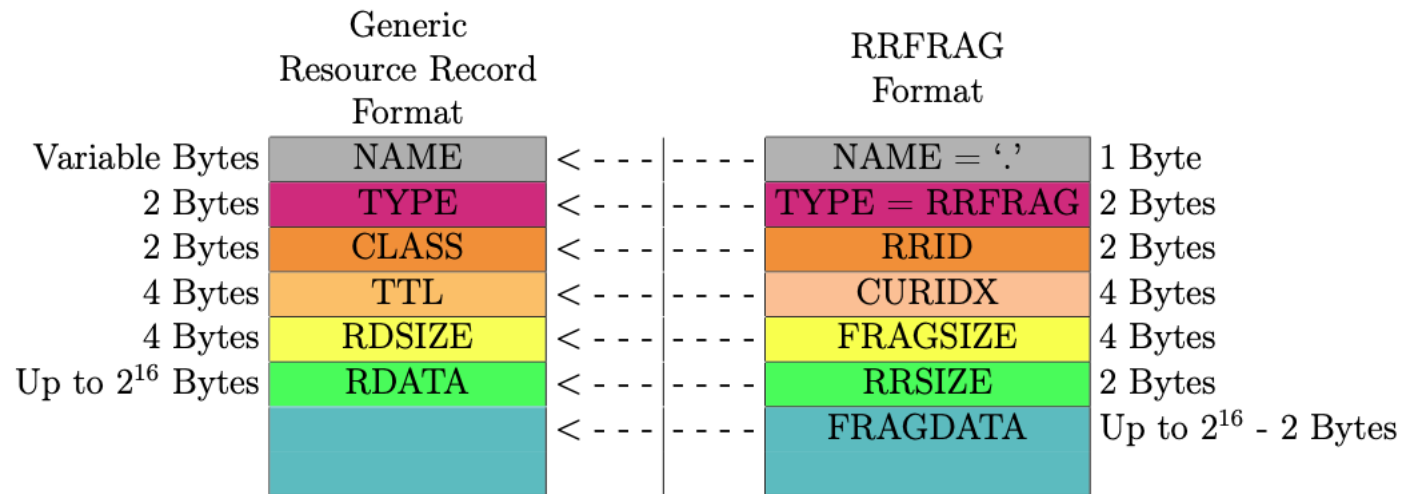
ARRF is a DNS-level fragmentation mechanism designed with backwards compatibility in mind

- Rather than breaking up a large DNS message, fragment the problem resource records
- Resolver makes explicit requests for the fragments it wants, so no advertising is required and prevents ICMP flooding
- Does not waste a round trip compared to falling back to TCP
- Can use parallelism to reduce total pain

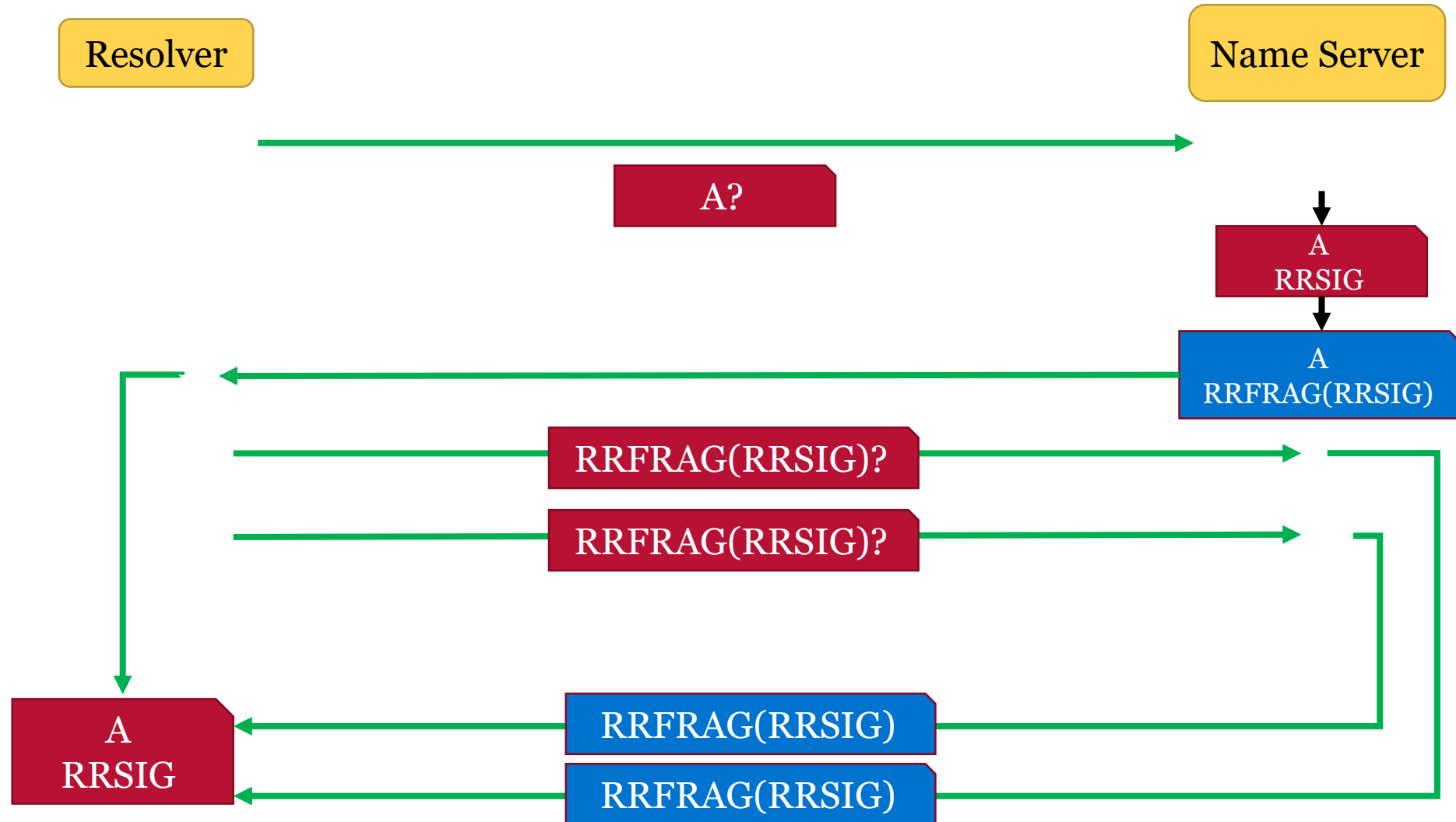
RRFRAG

When a resource record is too large, it is replaced with a Resource Record Fragment (RRFRAG)

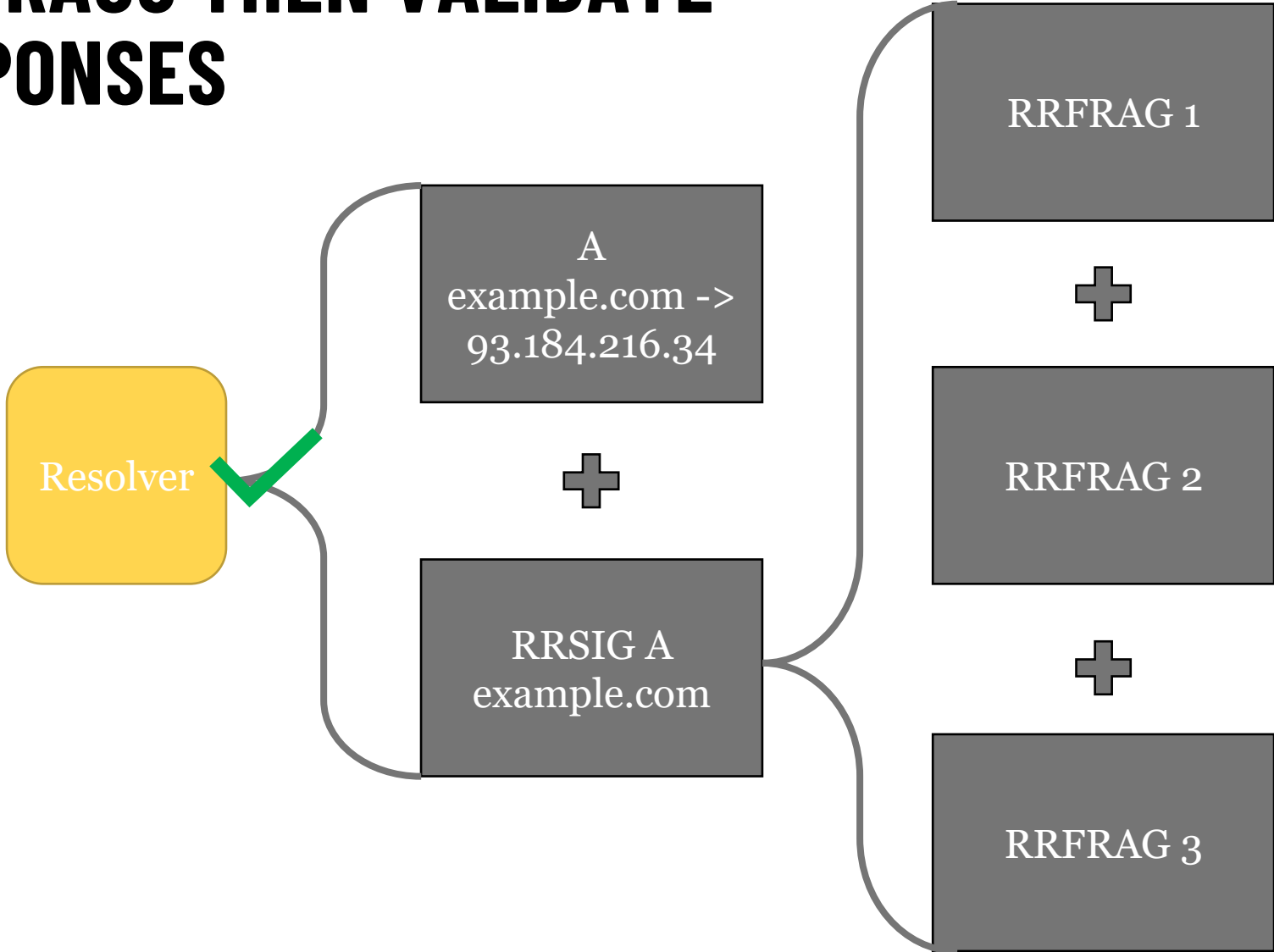
- Constructed to be mapped onto the Generic Resource Record Format to maximize compatibility



ARRF IN ACTION

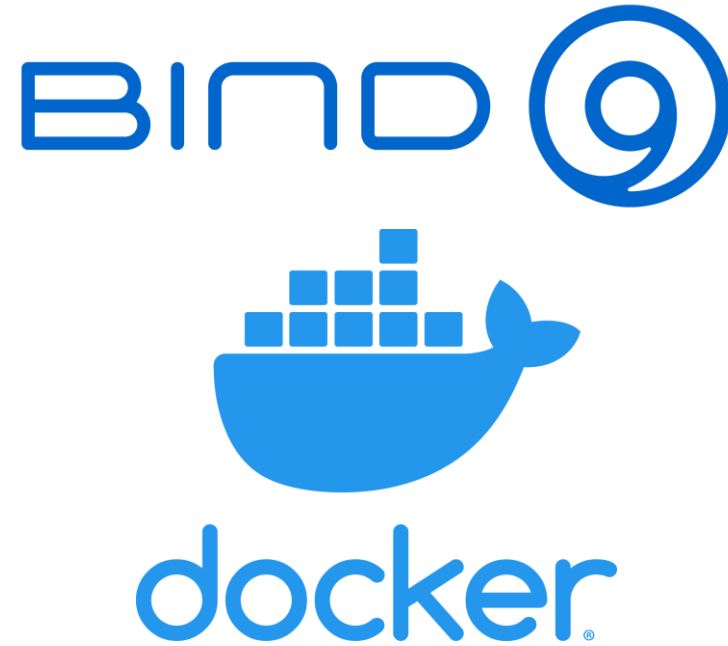


COMBINE RRFRAGS THEN VALIDATE DNSSEC RESPONSES



EXPERIMENTAL EVALUATION

- Used Open Quantum Safe's liboqs and OpenSSL fork to add Falcon-512, Dilithium2 and SPHINCS+-SHA256-128s support to Bind9
- Implemented a daemon which runs in front of Bind9 to transparently implement ARRF
- Construct a test DNS network using Docker containers running both Bind9 and ARRF daemon



EVALUATING PERFORMANCE

- Perform 1,000 queries for a unique 'A' record for each of the three algorithms
- Vary network conditions artificially
 - No delay, no Bandwidth restrictions
 - 10 ms delay 128 KBps bandwidth
 - 10 ms delay 50 MBps bandwidth
 - 100 ms delay 50 MBps bandwidth
- Vary the maximum DNS message size
- Evaluate the algorithms using Standard DNS (with TCP Fallback), sequential ARRF, and a parallelized version of ARRF

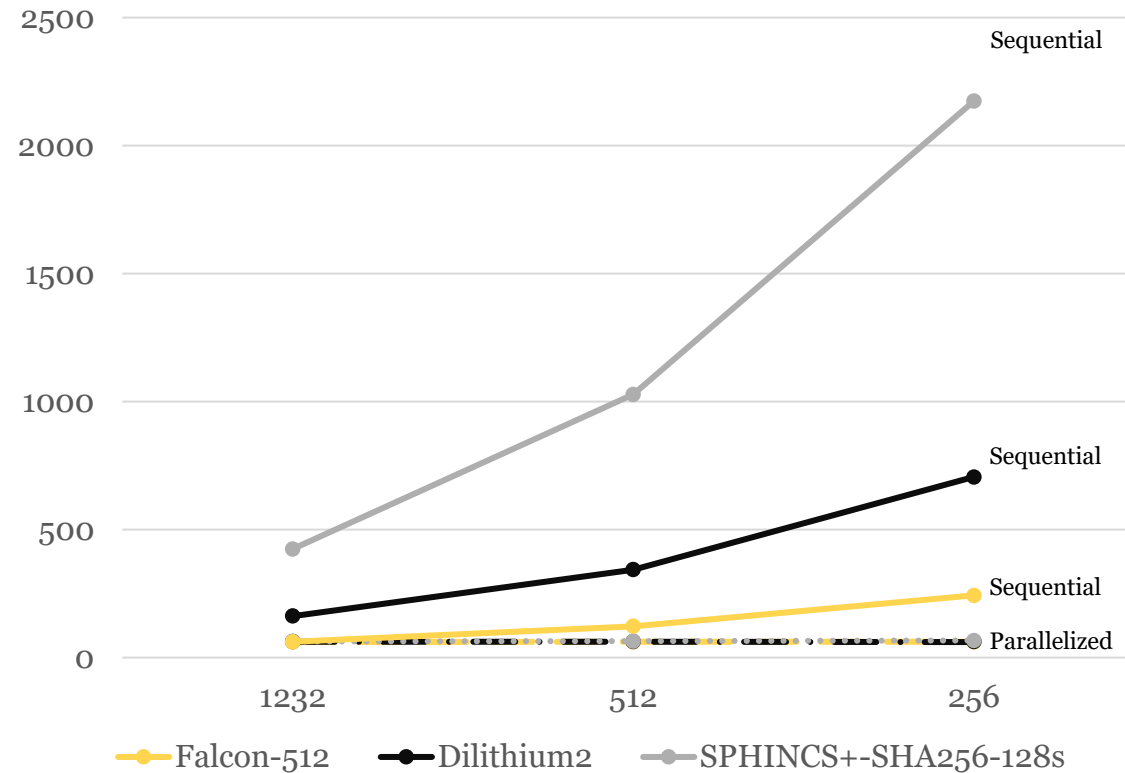
COMPARING LATENCY OF ARRF WITH STANDARD DNS

Resolution times (ms) with 10ms of latency and 50 Megabytes per second bandwidth
with a maximum size of 1232 bytes

Algorithms	Standard DNS with TCP fallback	Parallelized ARRF	Sequential ARRF
Falcon-512	82.11	61.96	62.07
Dilithium2	82.24	62.52	162.9
SPHINCS+-SHA256-128S	82.59	63.45	424.7
RSA 2048 with SHA256	41.50	-	-
ECDSA P256	47.49	-	-

PARALLELIZED ARRF LATENCY SCALES WELL WITH SMALLER PACKET SIZES

Resolution times (ms) with 10ms of latency and 50
Megabytes per second bandwidth for maximum DNS
message sizes 1232, 512, 256



DATA OVERHEAD OF ARRF COMPARED TO STANDARD DNS

Total bytes transmitted between resolver and name server during DNS lookup

Algorithm	Standard DNS with TCP fallback	ARRF		
		1232	512	256
Falcon-512	3,112	2,557	2,947	3,637
Dilithium2	8,623	8,367	9,402	11,322
SPHINCS+-SHA256-128S	26,073	26,140	29,620	36,175

ADDRESSING SECURITY CONCERNS

- ARRF does not make DNS based DoS attacks worse
 - All RRFAGs must be requested, and if a resolver receives one it is not expecting the RRFAG should be discarded
 - If an RRFAGs is modified while DNSSEC is being used, a validation failure will occur, which is no worse than a middle box modifying an RRSIG
- DNS Cache Poisoning is not a concern
 - RRFAG are never to be cached, thus not opening an avenue of attack for cache poisoning
- Memory Exhaustion Attacks
 - ARRF as specified is susceptible to Memory exhaustion attacks due to RRFAGs not having their integrity ensured

FUTURE WORK

- Addressing Memory exhaustion attacks
- ARRF was designed for backwards compatibility in mind, but testing on the real internet is needed
- How does ARRF perform on an unreliable network? Would likely need a timeout of some sort
- Does request-based fragmentation work for other protocols?
 - We think so! TurboTLS: <https://arxiv.org/abs/2302.05311>

Post-Quantum Signatures in DNSSEC via Request-Based Fragmentation

Jason Goertzen and Douglas Stebila

- ARRF is a lightweight, performant, and easy to implement modification to DNS protocol
- ARRF enables transmission of larger post-quantum signatures and keys in DNSSEC
- Removes the need for TCP fallback to send large DNS messages
- Performs 20% better than TCP when using parallelized ARRF
- Uses less data than TCP in cases of light-moderate fragmentation